CS 70 Discrete Mathematics and Probability Theory Summer 2019 James Hulett and Elizabeth Yang DIS 1D

1 Graph Basics

In the first few parts, you will be answering questions on the following graph G.



- (a) What are the vertex and edge sets V and E for graph G?
- (b) Which vertex has the highest in-degree? Which vertex has the lowest in-degree? Which vertices have the same in-degree and out-degree?
- (c) What are the paths from vertex *B* to *F*, assuming no vertex is visited twice? Which one is the shortest path?
- (d) Which of the following are cycles in G?
 - i. (B,C), (C,D), (D,B)
 - ii. (F,G), (G,F)
 - iii. (A,B), (B,C), (C,D), (D,B)
 - iv. (B,C), (C,D), (D,H), (H,G), (G,F), (F,E), (E,D), (D,B)
- (e) Which of the following are walks in *G*?

i. (E,G)ii. (E,G), (G,F)iii. (F,G), (G,F)iv. (A,B), (B,C), (C,D), (H,G)v. (E,G), (G,F), (F,G), (G,C)vi. (E,D), (D,B), (B,E), (E,D), (D,H), (H,G), (G,F)

(f) Which of the following are tours in G?

i. (E,G)ii. (E,G), (G,F)iii. (F,G), (G,F)iv. (E,D), (D,B), (B,E), (E,D), (D,H), (H,G), (G,F)

In the following three parts, let's consider a general undirected graph G with n vertices $(n \ge 3)$.

- (g) True/False: If each vertex of G has degree at most 1, then G does not have a cycle.
- (h) True/False: If each vertex of G has degree at least 2, then G has a cycle.
- (i) True/False: If each vertex of G has degree at most 2, then G is not connected.
- 2 Eulerian Tour and Eulerian Walk



- (a) Is there an Eulerian tour in the graph above?
- (b) Is there an Eulerian walk in the graph above?
- (c) What is the condition that there is an Eulerian walk in an undirected graph? Briefly justify your answer.

3 Odd Degree Vertices

Claim: Let G = (V, E) be an undirected graph. The number of vertices of G that have odd degree is even.

Prove the claim above using:

- (i) Direct proof (e.g., counting the number of edges in *G*). *Hint: in lecture, we proved that* $\sum_{v \in V} \deg v = 2|E|$.
- (ii) Induction on m = |E| (number of edges)
- (iii) Induction on n = |V| (number of vertices)

4 Trees

Recall that a *tree* is a connected acyclic graph (graph without cycles). In the note, we presented a few other definitions of a tree, and in this problem, we will prove two fundamental properties of a tree, and derive two definitions of a tree we learned from the note based on these properties. Let's start with the properties:

(a) Prove that any pair of vertices in a tree are connected by exactly one (simple) path.

(b) Prove that adding any edge (not already in the graph) between two vertices of a tree creates a simple cycle.

Now you will show that if a graph satisfies this property then it must be a tree:

(c) Prove that if the graph has no simple cycles and has the property that the addition of any single edge (not already in the graph) will create a simple cycle, then the graph is a tree.